# Nonlinear Laplacian for Digraphs and Its Applications to Network Analysis

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Can we develop spectral graph theory for digraphs?

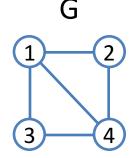
- Spectral graph theory analyzes graph properties via eigenpairs of associated matrices.
  - Adjacency matrix, incidence matrix, Laplacian
- Applications
  - Approximation to graph parameters (e.g, chromatic number), community detection, visualization, etc.
- Well established for undirected graphs.

Can we develop spectral graph theory for digraphs?

- Many real-world networks are directed!
  - Web graph, Twitter followers, phone calls, paper citations, food web, metabolic network.
- Extensions for digraphs are largely unexplored and unsatisfying.

# Laplacian

- Graph G = (V, E)
- Adjacency matrix: A<sub>G</sub>
- Degree matrix: D<sub>G</sub>
- Laplacian  $L_G := D_G A_G$

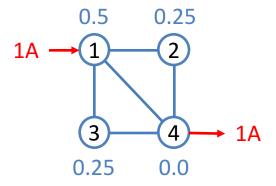


• Normalized Laplacian  $\mathcal{L}_{G} := D_{G}^{-1/2} L_{G} D_{G}^{-1/2} = I - D_{G}^{-1/2} A_{G} D_{G}^{-1/2}$ 

### **Interpretation of Laplacian**

- Regard G as an electric circuit.
- An edge = a resistance of  $1\Omega$ .
- Flow a current of  $\mathbf{b}(\mathbf{u})$  ampere to each vertex  $\mathbf{u} \in \mathbf{V}$ .

The voltages of vertices can be computed by solving  $L_G \mathbf{x} = \mathbf{b}$ 



Existing extensions of Laplacians for digraphs:

- 1.  $L_G = D_G^+ A_G$ 
  - Asymmetric and hence eigenpairs are complex-valued.
- 2. Chung's Laplacian
  - Assume strong connectivity. Need random walks to interpret its eigenpairs.

#### **Our contributions**

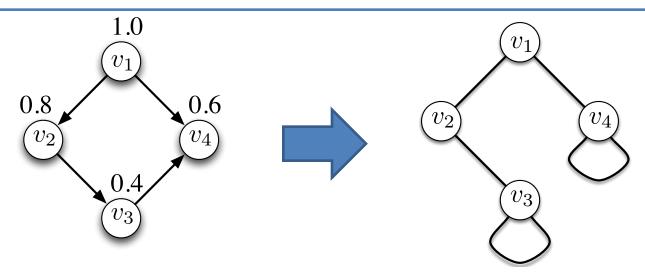
- 1. Laplacian for digraphs whose eigenpairs can be interpreted more combinatorially.
- 2. Algorithm that computes a small eigenvalue.
- 3. Applications to visualization and community detection.

# **Nonlinear Laplacian**

Nonlinear Laplacian  $L_G: \mathbb{R}^n \rightarrow \mathbb{R}^n$  for a digraph G:

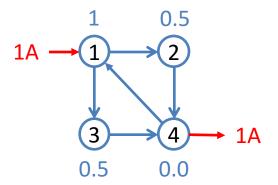
From a vector  $\mathbf{x} \in \mathbb{R}^n$ , we compute  $L_G(\mathbf{x})$  as follows

- 1. Define an *undirected* graph as follows: for each arc  $u \rightarrow v$ 
  - If  $\mathbf{x}(u) \ge \mathbf{x}(v)$ , add an (undirected) edge {u, v}.
  - Otherwise, add self-loops.
- 2. Let  $L_H$  be the Laplacian of H.
- 3. Output L<sub>H</sub>**x**.



- Regard G = (V, E) as an electric circuit.
- An edge = a diode of  $1\Omega$  (current flows only one way).
- Flow a current of  $\mathbf{b}(\mathbf{u})$  ampere to each vertex  $\mathbf{u} \in \mathbf{V}$ .

The voltages of vertices can be computed by solving  $L_G(\mathbf{x}) = \mathbf{b}$ .



# **Eigenpair of Nonlinear Laplacian**

- Normalized Laplacian  $\mathcal{L}_{G} : \mathbf{x} \mapsto D_{G}^{-1/2} L_{G} (D_{G}^{-1/2} \mathbf{x})$
- $(\lambda, \mathbf{v})$  is an eigenpair of  $\mathcal{L}_{G}$  if  $\mathcal{L}_{G}(\mathbf{v}) = \lambda \mathbf{v}$ - Trivial eigenpair:  $(\lambda_{1} = 0, \mathbf{v}_{1})$ .

For any subspace U of positive dimension,  $\Pi_U \mathcal{L}_G$  has an eigenpair. ( $\Pi_U$  = Projection matrix to U)

 $\Rightarrow$  Nontrivial eigenpair of  $\mathcal{L}_{G}$  exists by choosing U =  $\mathbf{v}_{1}^{\perp}$ .

Let  $\lambda_2$  be the smallest eigenvalue orthogonal to  $\mathbf{v}_1$ .

# Algorithm

Computing  $\lambda_2$  is (likely to be) NP-hard.

Suppose we start the diffusion process

$$d\boldsymbol{x} = -\Pi_U \mathcal{L}_G(\boldsymbol{x}) dt$$

from a vector in the subspace U =  $\mathbf{v}_1^{\perp}$ .

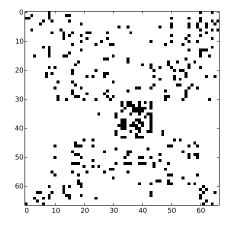
- **x** converges to an eigenvector orthogonal to  $\mathbf{v}_1$ .
- Rayleigh quotient  $\mathcal{R}_G(\mathbf{x}) := \frac{\mathbf{x}^T \Pi_U \mathcal{L}_G(\mathbf{x})}{\mathbf{x}^T \mathbf{x}}$  never increases during the process.
- $\Rightarrow$  We can get a eigenvector of a small eigenvalue.

Friendship network at a high school in Illinois

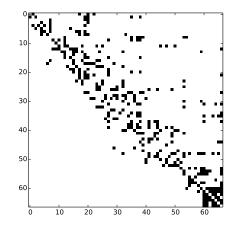
•  $u \rightarrow v$ : u regards v as a friend.

Reorder vertices according to the eigenvector computed by the diffusion process.

Chung's Laplacian



Nonlinear Laplacian



Our method shows the directivity of the network more clearly.

# **Visualization: Interpretation**

Laplacian for undirected graphs

$$\lambda_2 = \min \sum_{\{u,v\} \in E} (x(u) - x(v))^2$$

s.t. 
$$\|\mathbf{x}\| = 1$$
,  $\mathbf{x} \perp \mathbf{v}_1$ 

• Adjacent vertices are placed near.

#### Chung's Laplacian

$$\lambda_2 = \min \sum_{u \to v \in E} (\mathbf{x}(u) - \mathbf{x}(v))^2 \pi_u / d_u^+ \quad \text{s.t. } \|\mathbf{x}\| = 1, \mathbf{x} \perp \mathbf{v}_1.$$

Important vertices (w.r.t. RW) are placed in the middle.
<u>Nonlinear Laplacian</u>

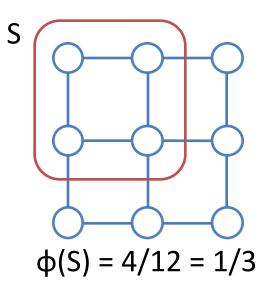
$$\lambda_2 = \min \sum_{u \to v \in E} \max(\mathbf{x}(u) - \mathbf{x}(v), 0)^2$$

s.t. 
$$\|\mathbf{x}\| = 1, \mathbf{x} \perp \mathbf{v}_1$$
.

- If  $\mathbf{x}(\mathbf{u}) \leq \mathbf{x}(\mathbf{v})$ , then we get no penalty.
- In particular,  $\lambda_2 = 0$  when G is a DAG.

### **Community Detection: Undirected Graphs**

S: Vertex set vol(S): Total degree of vertices in S cut(S): # of edges between S and V-S The conductance  $\phi(S)$  of S is  $\frac{\text{cut}(S)}{\min(\text{vol}(S), \text{vol}(V-S))}$ 



Small conductance

 $\rightarrow$  Good community

# **Community Detection: Undirected Graphs**

Cheeger's inequality ('70)  $\lambda_2/2 \le \min_S \varphi(S) \le \sqrt{(2\lambda_2)}$ 

- We can efficiently compute S with  $\phi(S) \leq v(2\lambda_2)$  from  $v_2$ .
- Still widely used.

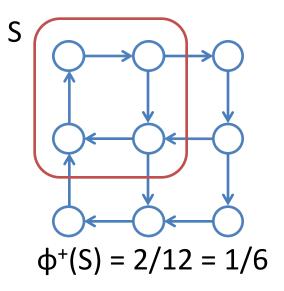
# **Community Detection: Digraphs**

S: Vertex set

vol(S): Total indegrees + outdegrees of vertices in S

cut<sup>+</sup>(S): # of arcs from S to V-S

(Directed) conductance  $\phi(S)$  of S is  $\frac{\min(\operatorname{cut}^+(S),\operatorname{cut}^+(V-S))}{\min(\operatorname{vol}(S),\operatorname{vol}(V-S))}$ 



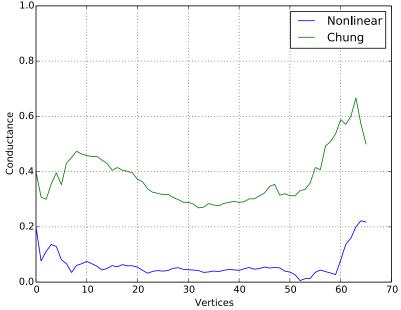
### **Community Detection: Digraphs**

Cheeger's inequality for digraphs

$$\lambda_2/2 \le \min_{S} \phi(S) \le 2\sqrt{\lambda_2}$$

• We can efficiently compute S with  $\phi(S) \leq 2\sqrt{\mathcal{R}_G}(\mathbf{x})$  from  $\mathbf{x}$ .

Reorder vertices according to the obtained eigenvector in the high school network, and plot  $\phi$  of each prefix set.



- φ is low everywhere = directivity
- φ rapidly increases = community

# Summary

Nonlinear Laplacian for digraphs

- Strong connectivity is not needed.
- Eigenpairs are combinatorially interpretable.
- Applications to visualization and community detection.

#### Future Work

- Approximation of  $\lambda_2$ .
- Finding a community in time proportional to its size.
- Other applications.