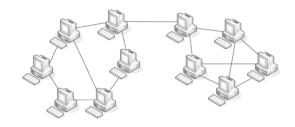
DL-Lite_R in the light of propositional logic for decentralized data management



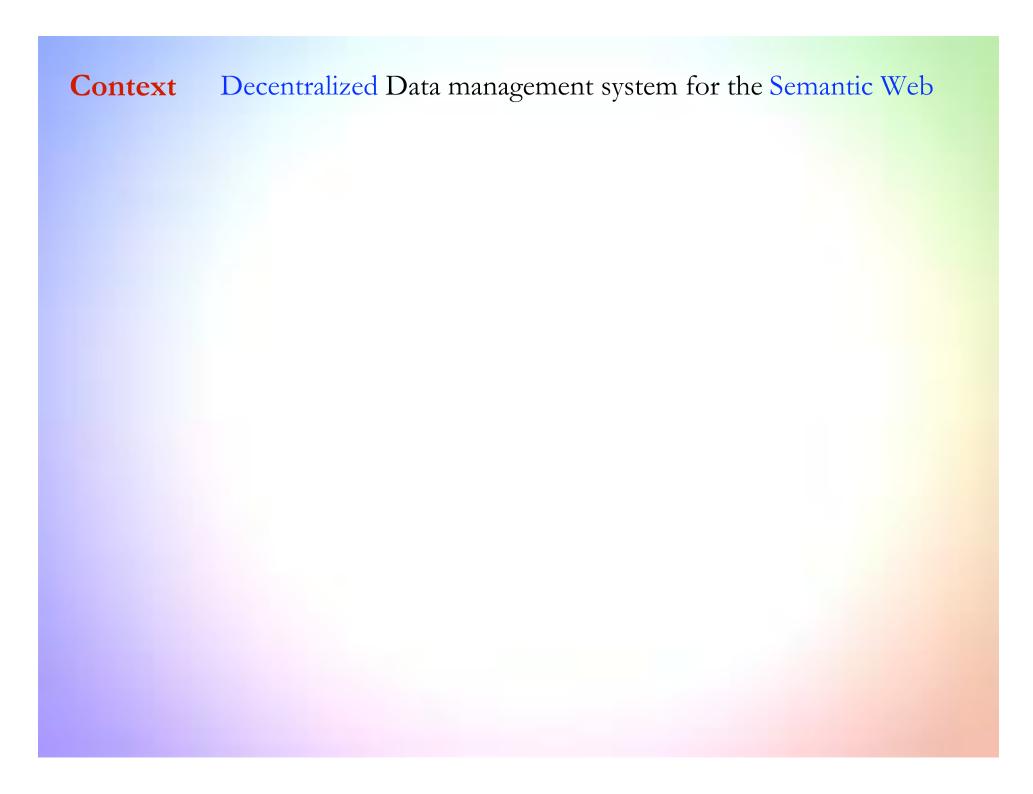
Nada Abdallah

François Goasdoué

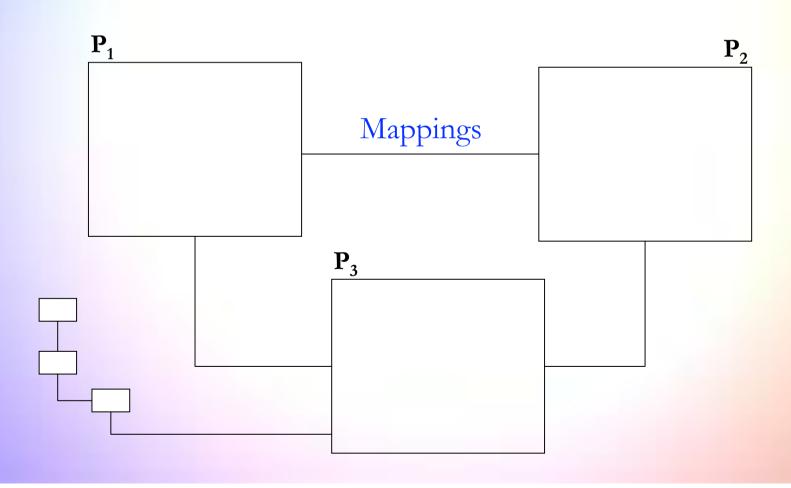
Université Paris-Sud & CNRS(LRI/IASI) – INRIA (Saclay/GEMO)

Marie-Christine Rousset

Université Grenoble & CNRS(LIG) – INRIA

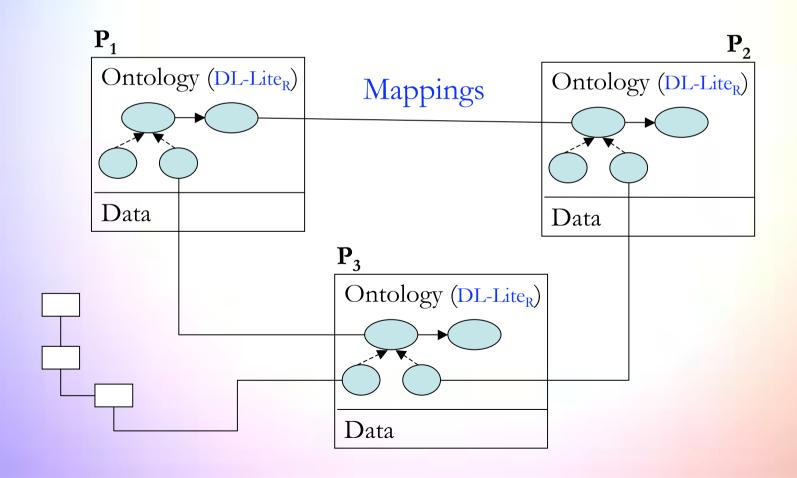


Decentralized: Dynamic Network of collaborative peers



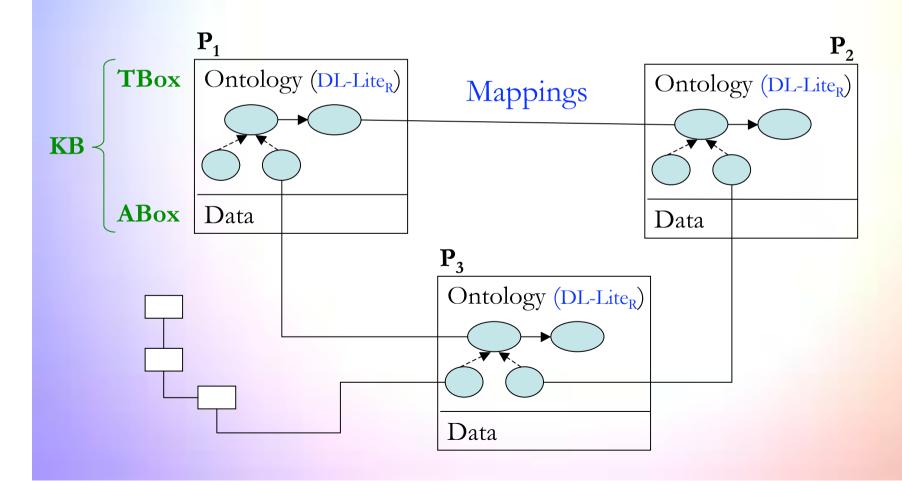
Decentralized: Dynamic Network of collaborative peers

Semantic Web: Data are described with ontologies (OWL2)



Decentralized: Dynamic Network of collaborative peers

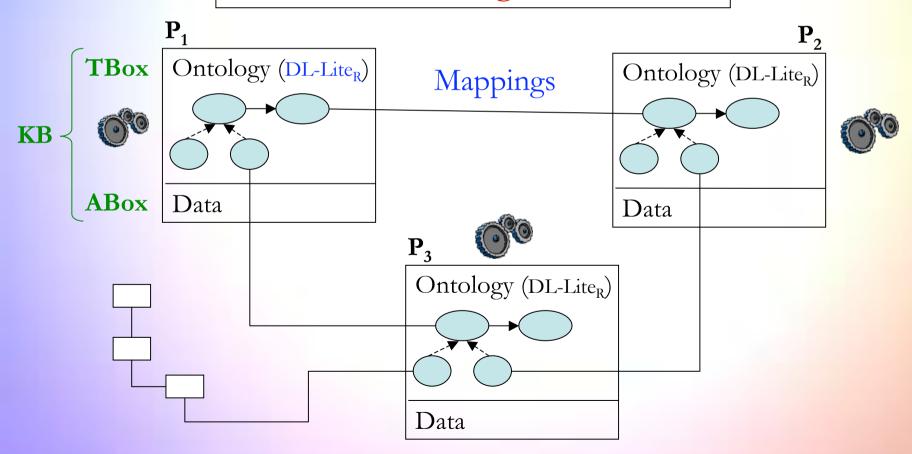
Semantic Web: Data are described with ontologies (OWL2)



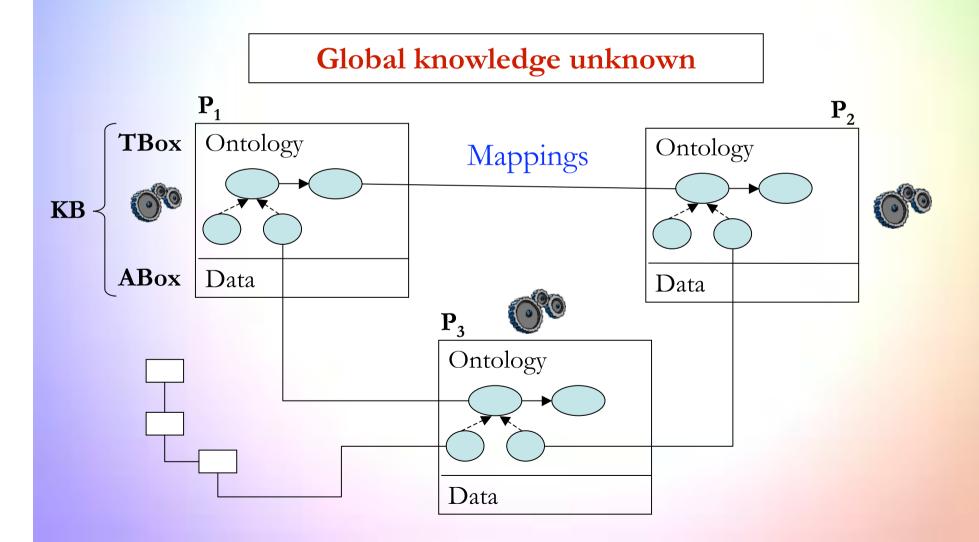
Decentralized: Dynamic Network of collaborative peers

Semantic Web: Data are described with ontologies (OWL2)

Global knowledge unknown

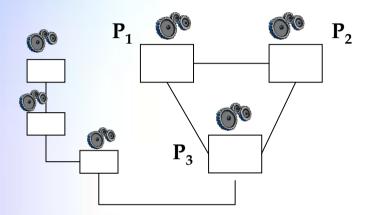


EXISTING: SomeOWL - SomeRDFS



Contributions

1. DL-Lite_R decentralized data model

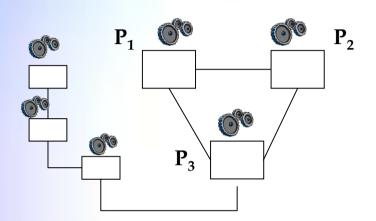


- 2. Decentralized Algorithms for:
 - ✓ Data consistency checking
 - ✓ Query Answering

Contributions

State of the art

1. DL-Lite_R decentralized data model



- 2. Decentralized Algorithms for:
 - ✓ Data consistency checking
 - **✓** Query Answering

• D.Calvenese and al. (JAR 07)

"Tractable reasoning and efficient query answering in description logics"

Centralized DL-Lite_R KB

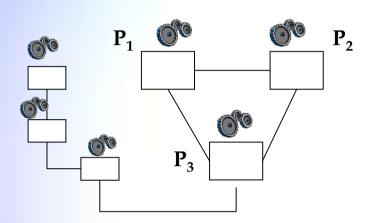
Centralized Algorithms for:

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Contributions

State of the art

1. DL-Lite_R decentralized data model



- 2. Decentralized Algorithms for:
 - ✓ Data consistency checking
 - **✓** Query Answering

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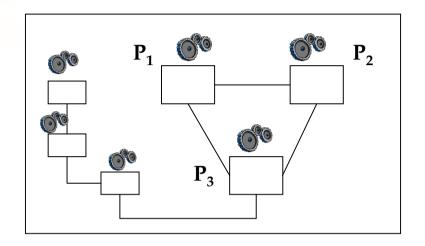
"Tractable reasoning and efficient query answering in description logics"

Centralized DL-Lite_R KB

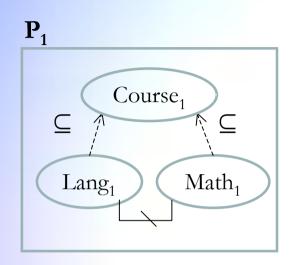
Centralized Algorithms for:

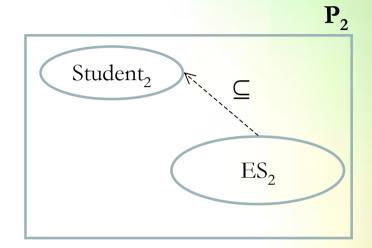
- ✓ Data consistency checking
- ✓ Query Answering
- 3. Decentralized & Centralized Algorithms for:
 - ✓ View consistency checking
 - ✓ Query Answering using views

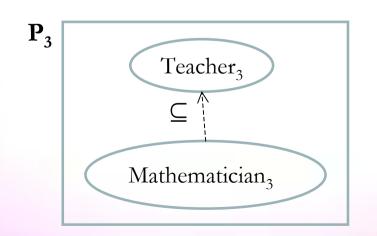
Outline

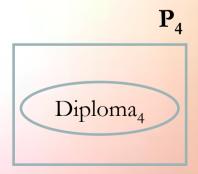


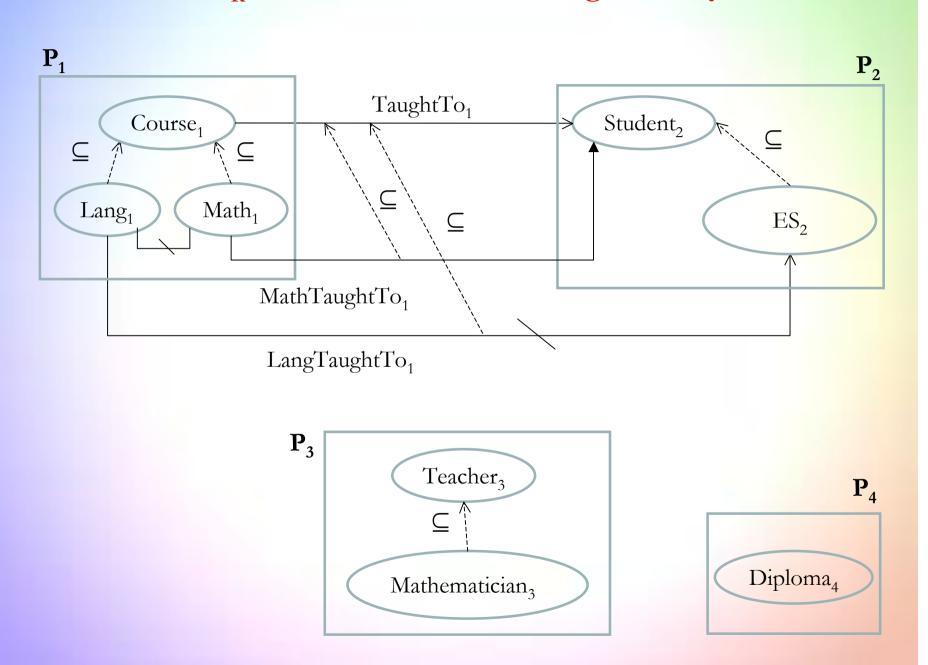
- ☐ DL-Lite_R decentralized data management system
 - ☐ Data consistency checking
 - ☐ Query Answering
 - ☐ Conclusion & Future Work

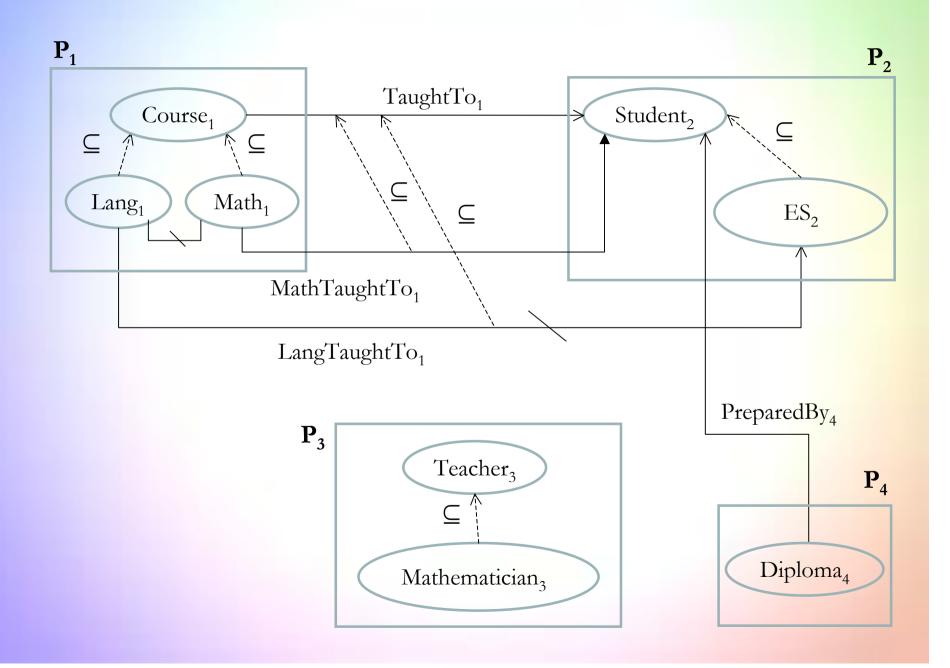


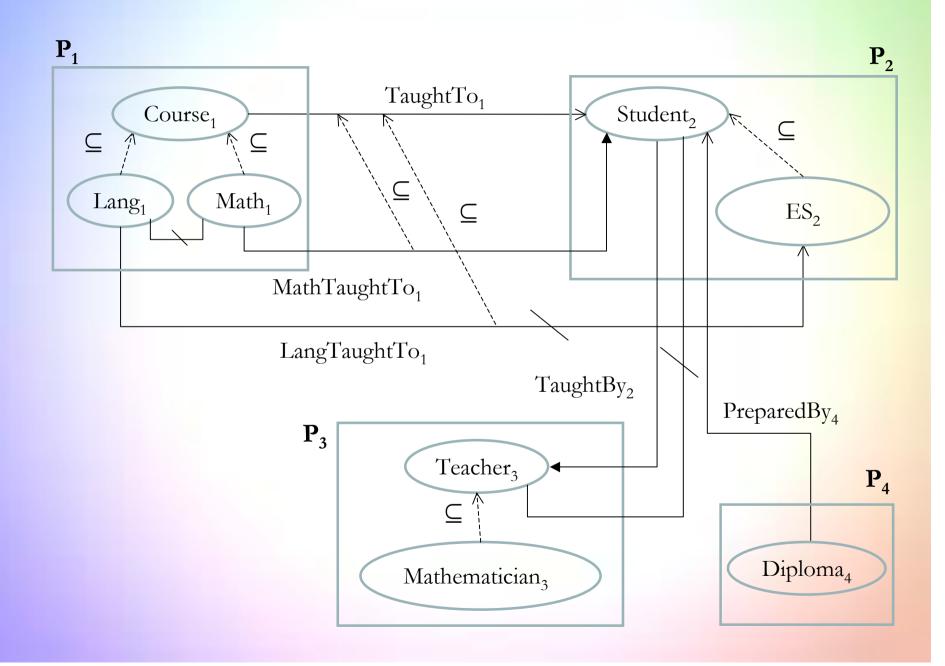


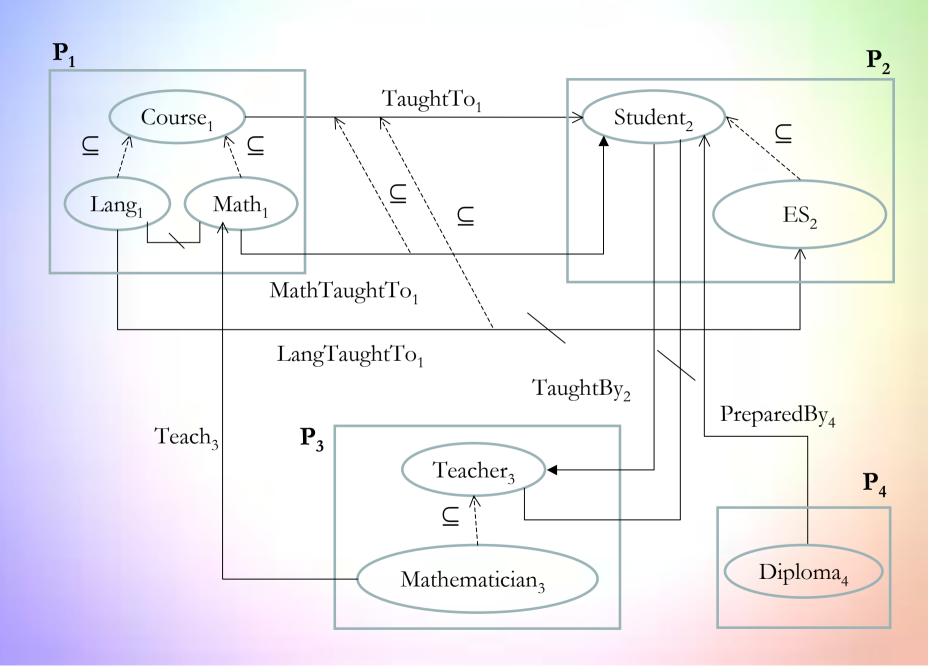


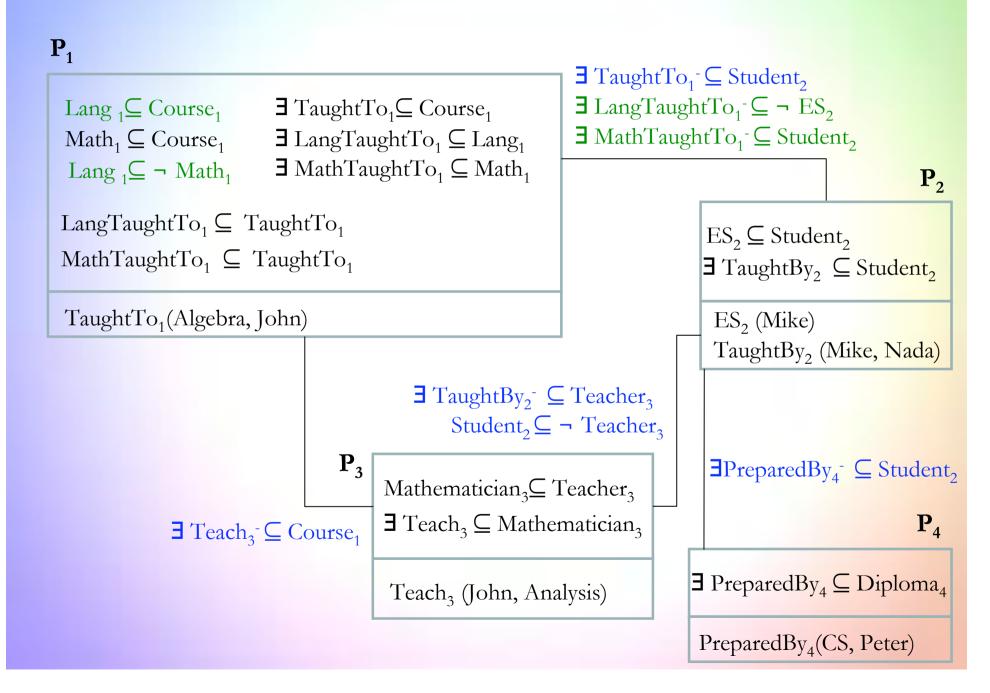




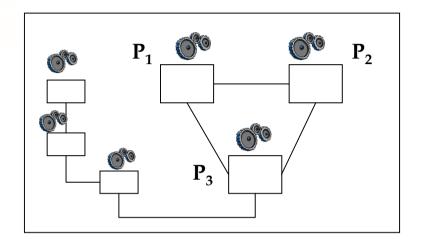








Outline



- ☐ DL-Lite_R decentralized data management system
 - ☐ Data consistency checking
 - ☐ Query Answering
 - ☐ Conclusion & Future Work

• Consistency problem:

A data management system is consistent iff its KB (Tbox + Abox) is satisfiable.

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- In DL-Lite_R:
 - ✓ Tbox is always satisfiable
 - ✓ Tbox + Abox may be insatisfiable

Consistency problem:

A data management system is consistent iff its KB (Tbox + Abox) is satisfiable.

• In DL-Lite_R:

- ✓ Tbox is always satisfiable
- ✓ Tbox + Abox may be insatisfiable

• Example:

TBox \exists Taught To_1 \subseteq Student₂ Student₂ \subseteq \neg Teacher₃ Mathematician₃ \subseteq Teacher₃ \exists Teach₃ \subseteq Mathematician₃

ABox TaughtTo₁ (Algebra, John)
Teach₃ (John, Analysis)

Implicit Negative Inclusion

TBox $\models \exists TaughtTo_1 \subseteq \neg \exists Teach_3$ KB $\models \exists y TaughtTo_1 (y, John)$ KB $\models \exists y Teach_3 (John, y)$

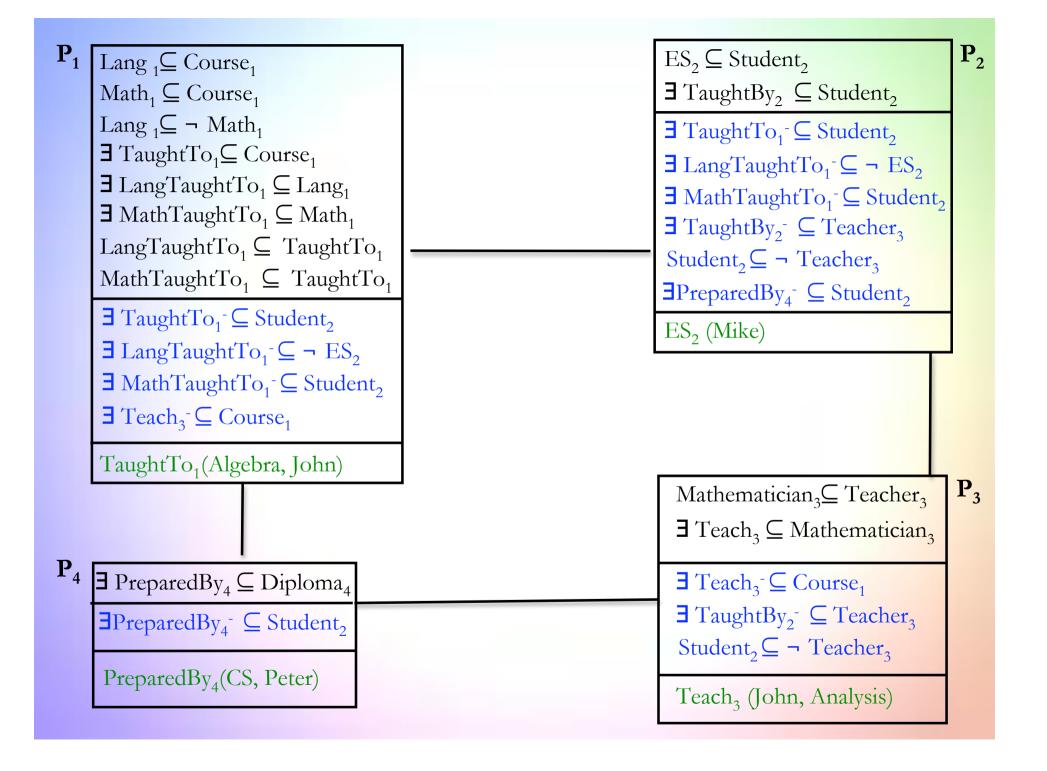
Centralized case

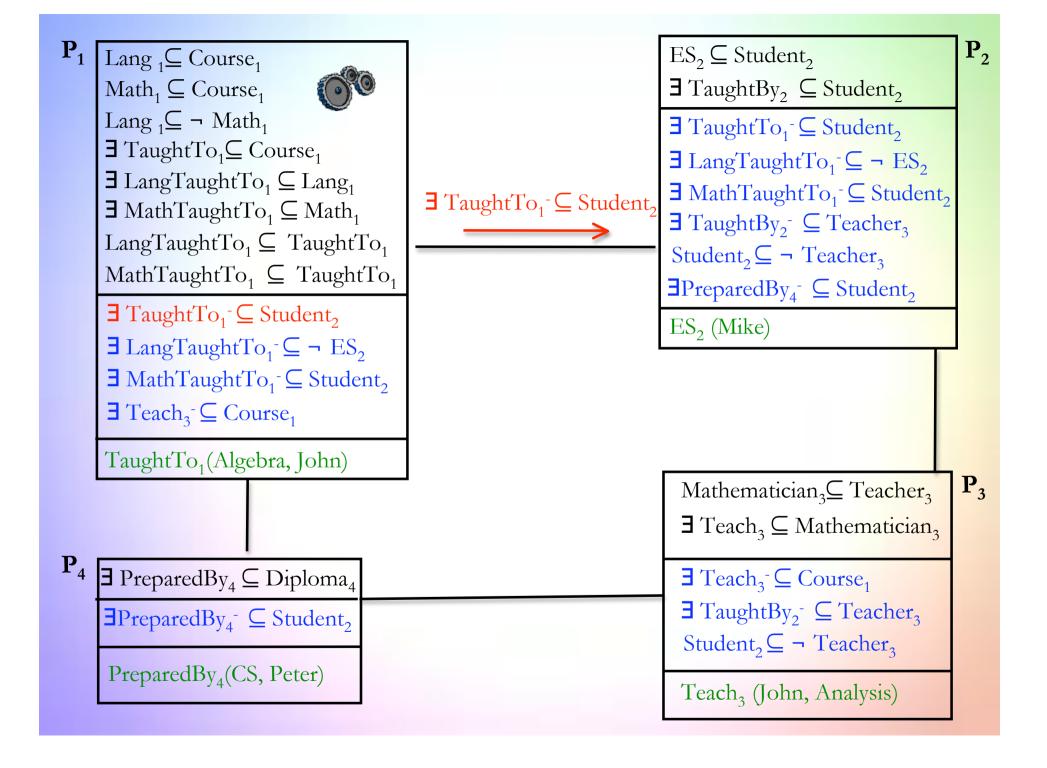
Existing Algorithm Consistent:

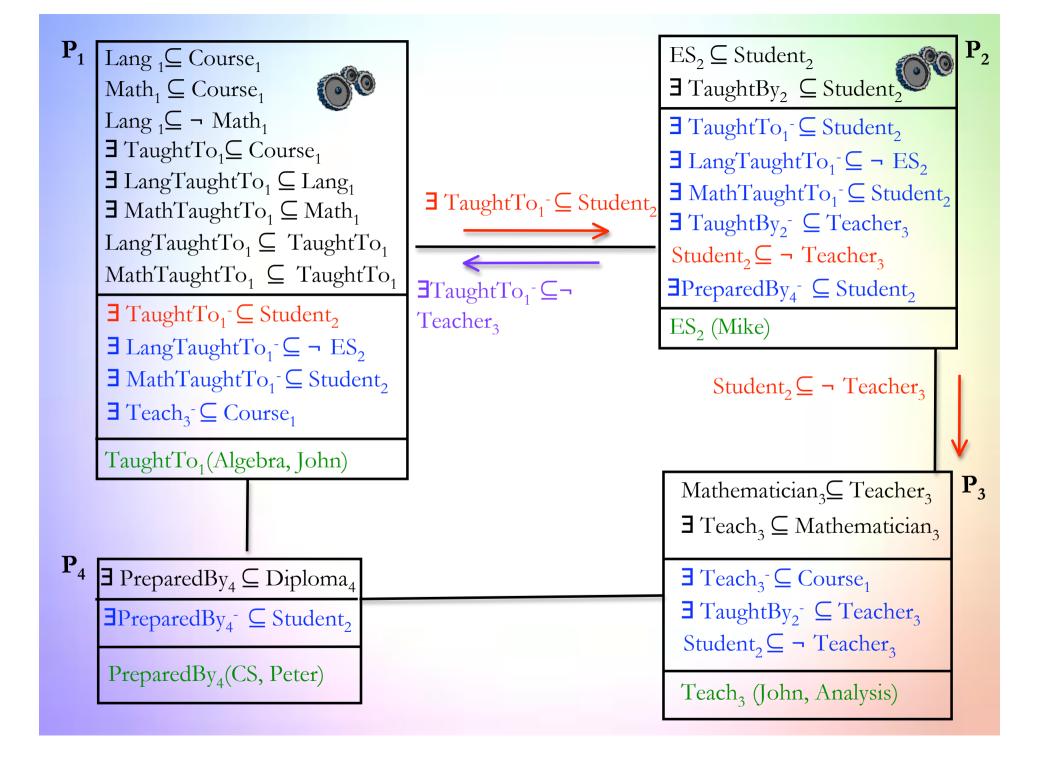
- D.Calvenese and al. (JAR 07): "Tractable reasoning and efficient query answering in description logics"
- Computes all the Negative Inclusions entailed by the Tbox.
- > Checks if the computed negative inclusions are violated by the Abox.

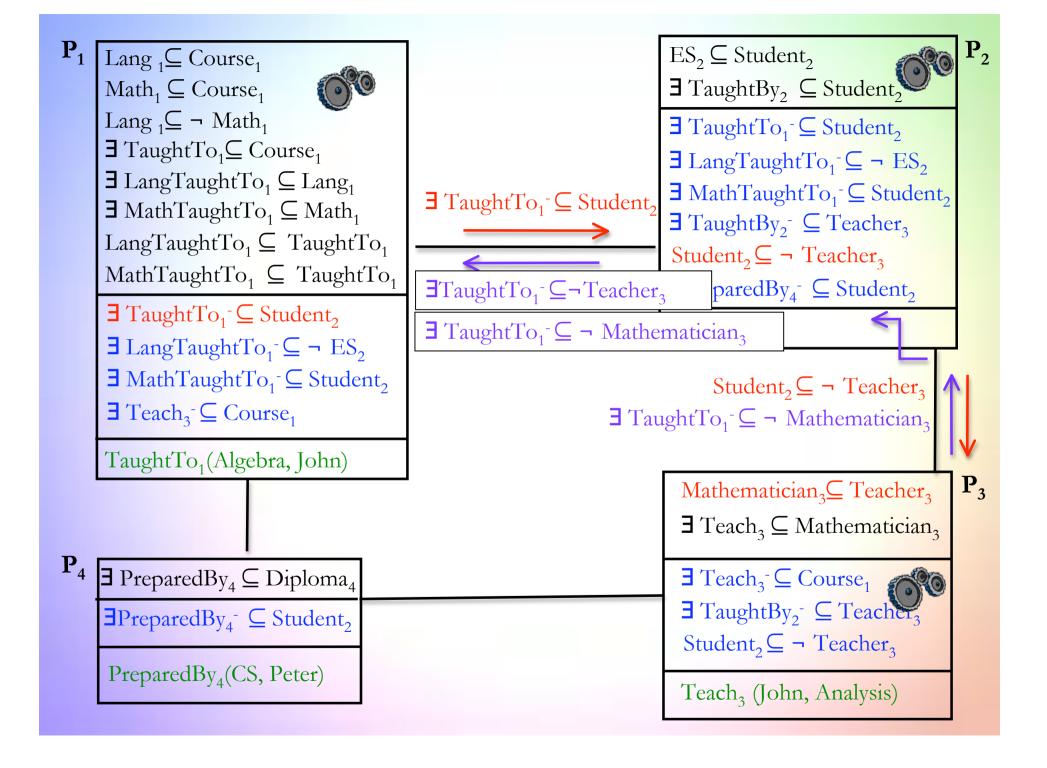
Checking if the DL-Lite_R formulae that must be disjoint according to the TBox, indeed have disjoint instances in the Abox

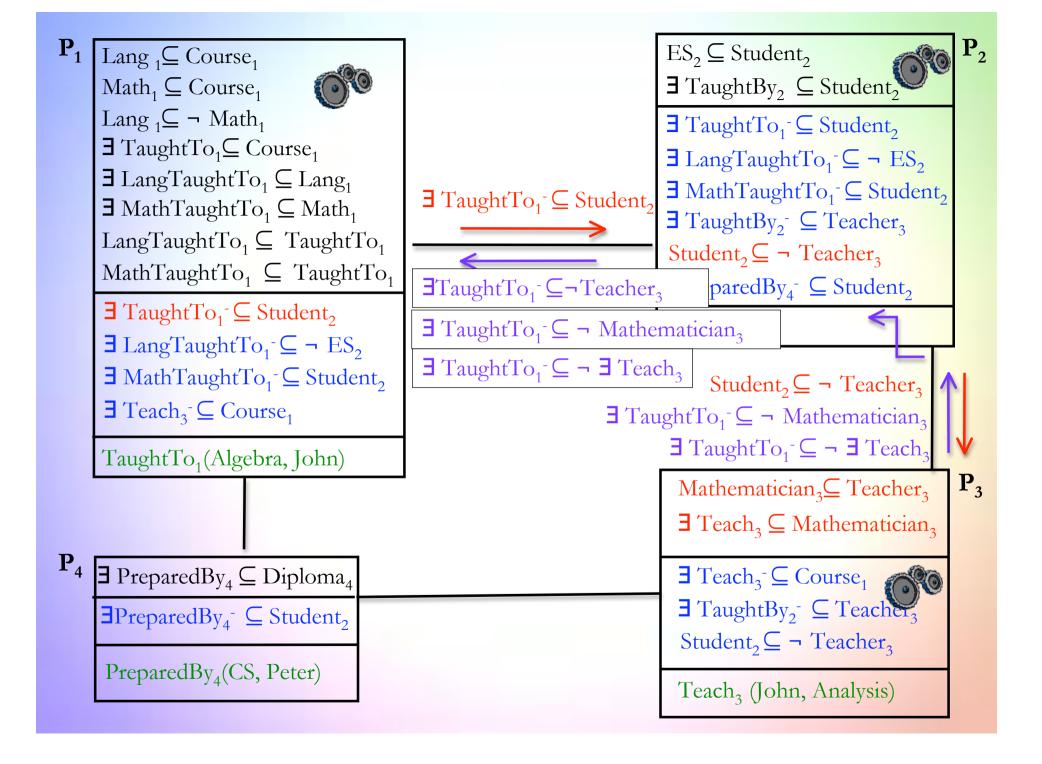
- Decentralized case
 - Decentralized Data consistency Checking w.r.t a peer
 - ✓ Each peer propagates over the network each of its positive or negative inclusions and collects in return all the negative inclusions of the network the entailment of which uses the Tbox of the peer.
 - o All Negative Inclusions entailed by the network are not necessarily computed.
 - ✓ Each peer checks if the negative inclusions that he has collected are violated by the Data by sending queries to the corresponding peers.

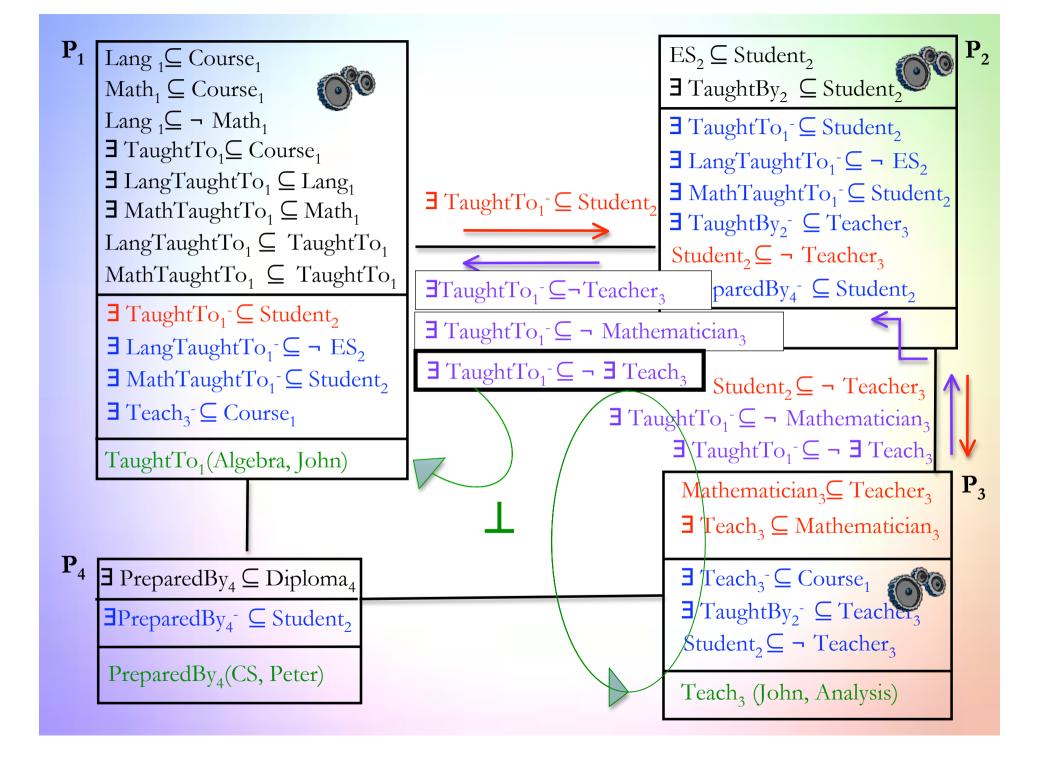


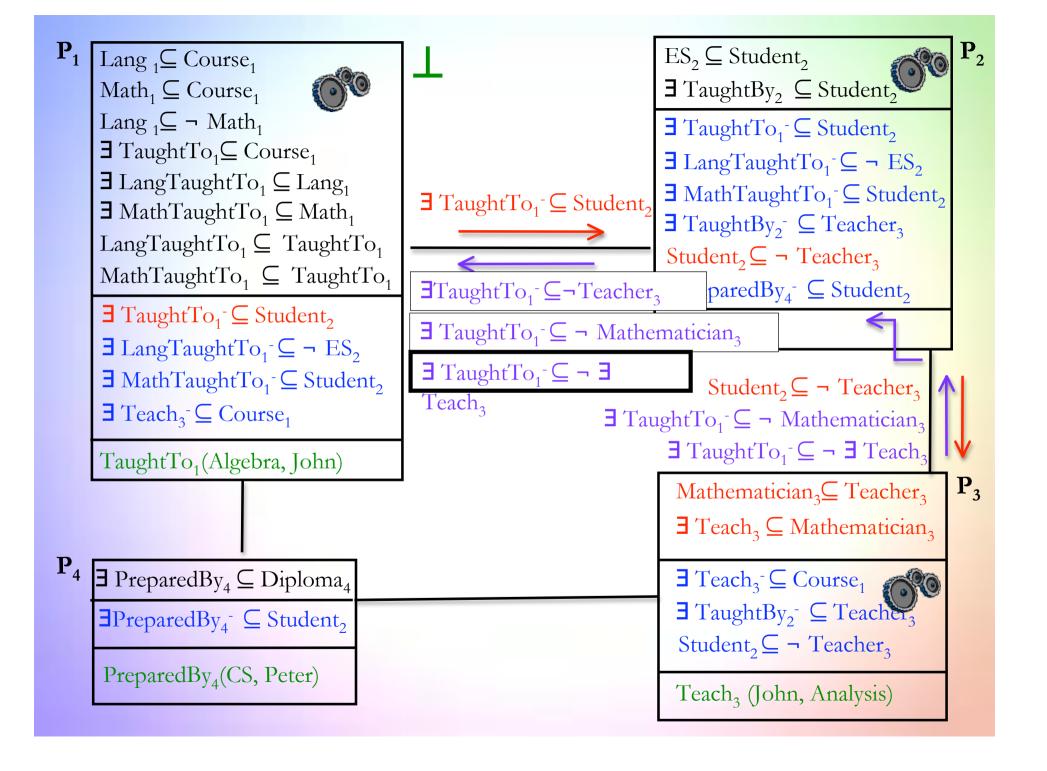


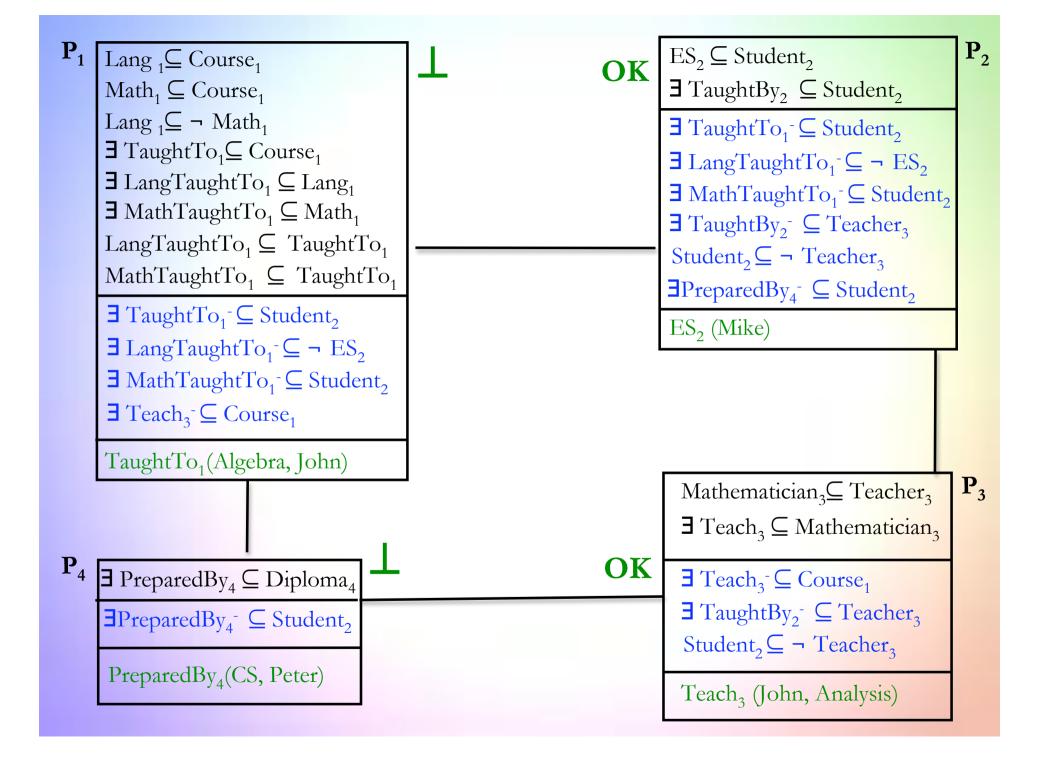






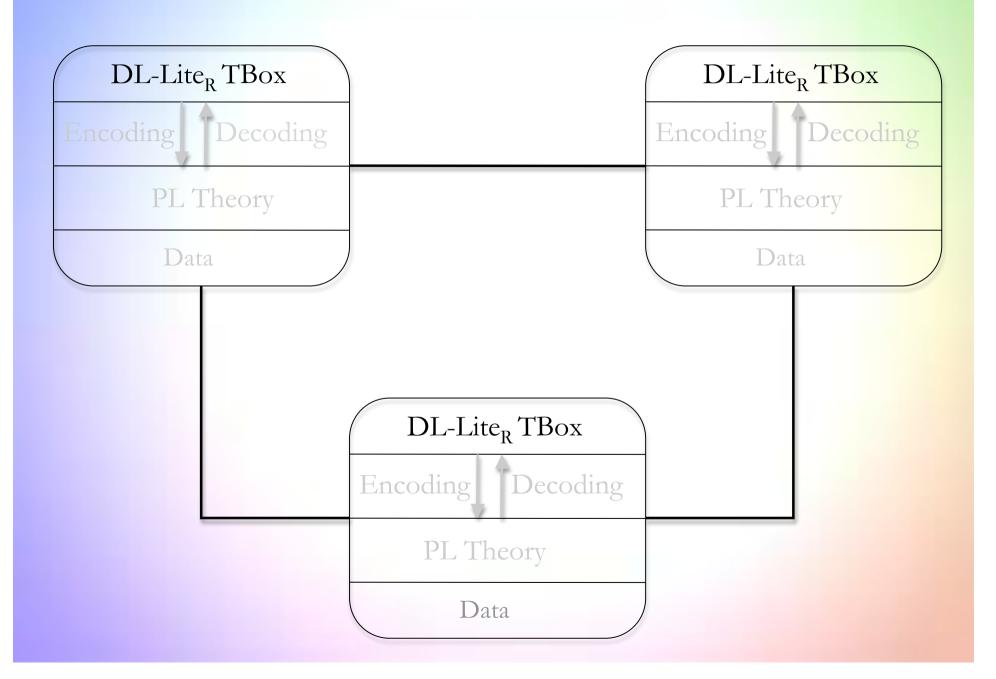


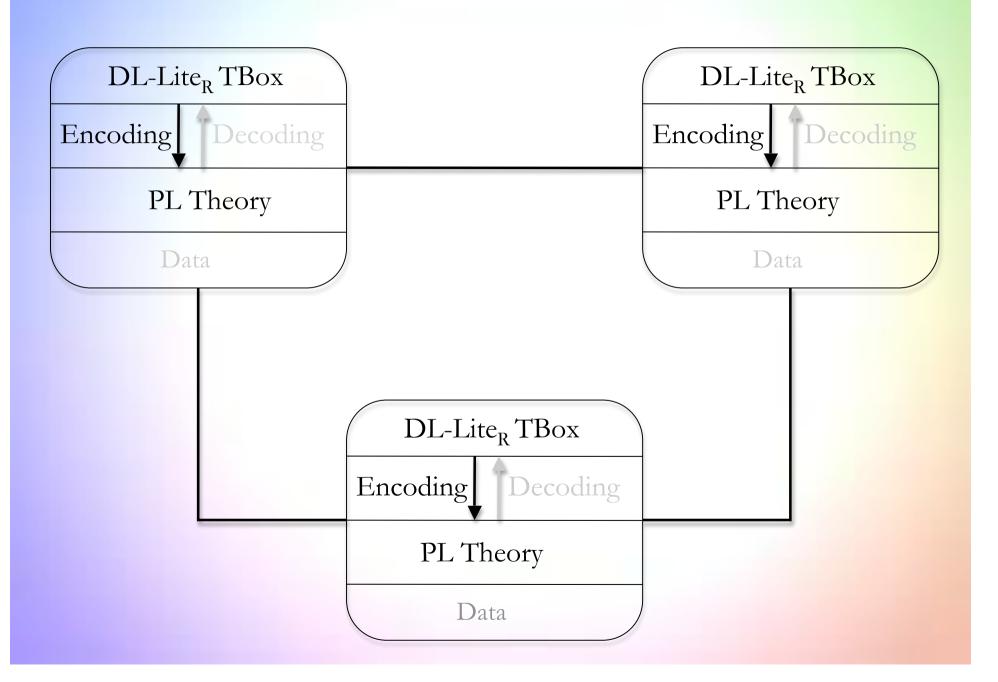


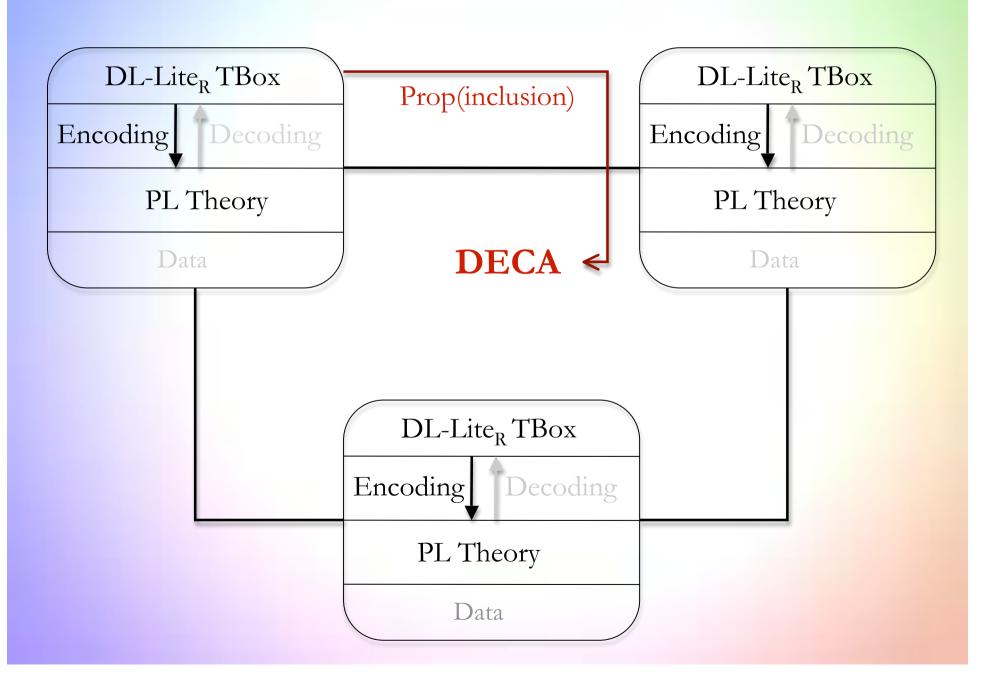


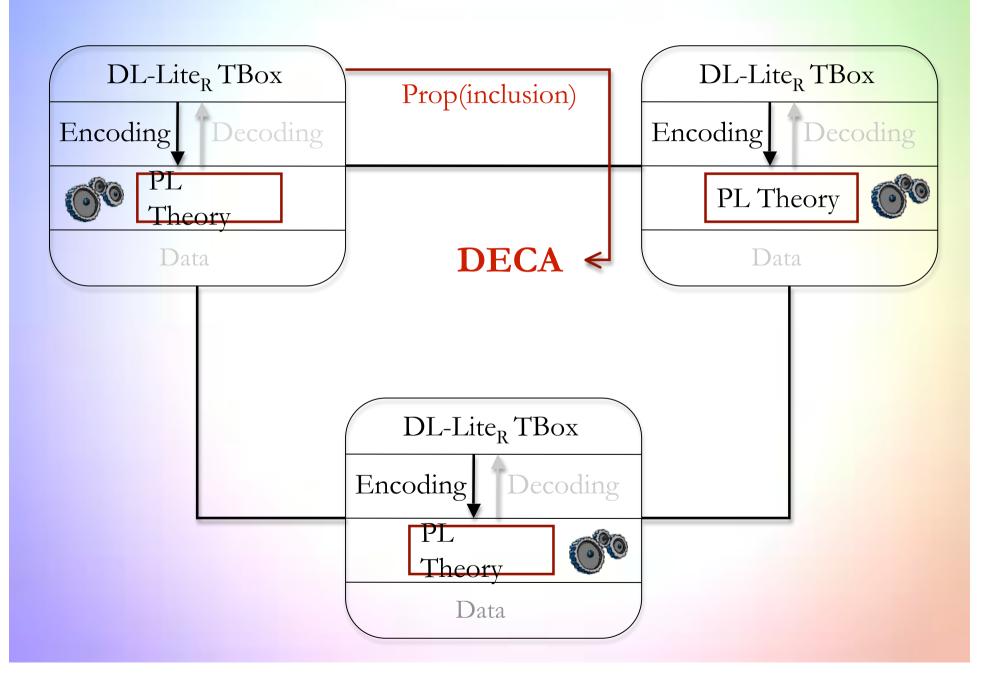
Decentralized Negative Inclusion computation by using SomeWhere

- ✓ An existing Decentralized Infrastructure "SomeWhere" for reasoning in propositional logic and for which the experiments have demonstrated the scalability.
 - P. Adjiman and al. (IJCAI 05): Scalability study of peer-to-peer consequence finding"
- ✓ An existing Decentralized Consequence finding Algorithm in propositional logic (DECA)
- ✓ A propositional encoding of the DL-Lite_R TBox.

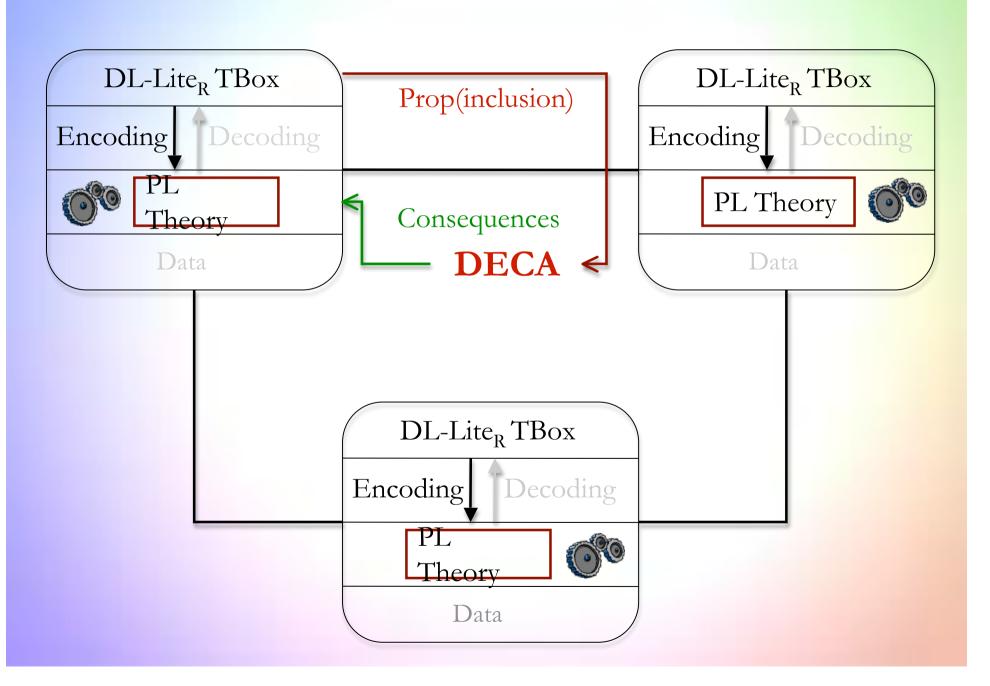






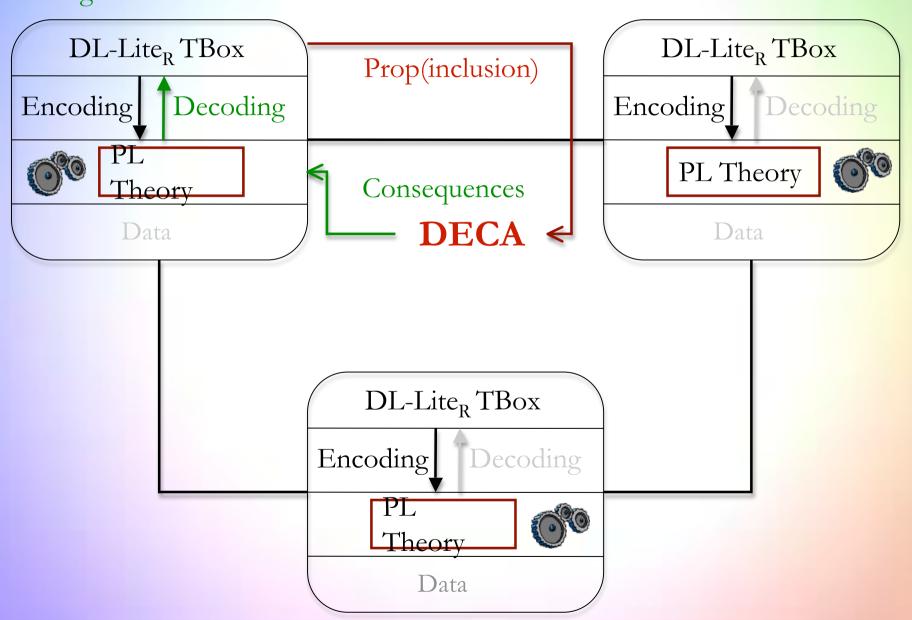


Implementation using SomeWhere



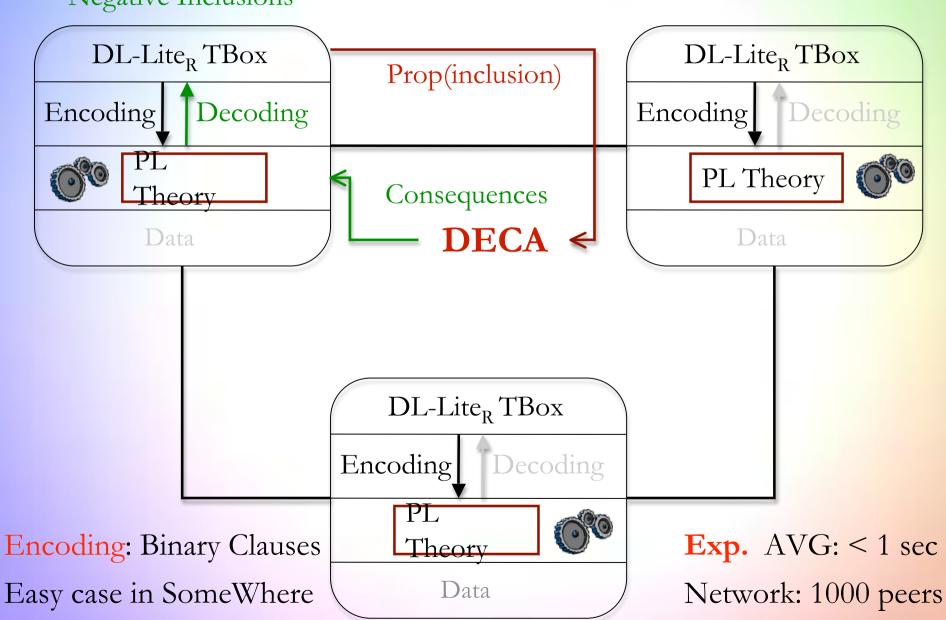
Implementation using SomeWhere

Negative Inclusions

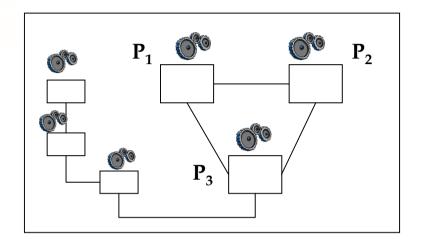


Implementation using SomeWhere

Negative Inclusions



Outline



- ☐ DL-Lite_R decentralized data management system
 - ☐ Data consistency checking
 - ☐ Query Answering
 - ☐ Conclusion & Future Work

Query Answering

• Query Answering problem:

Given a conjunctive query q, compute the answers of q over the system.

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• Answer:

A tuple t of the constants appearing in KB such that:

$$KB \models q(t)$$

Query Answering

• Query Answering problem:

Given a conjunctive query q, compute the answers of q over the system.

• Answer:

A tuple t of the constants appearing in KB such that: $KB \models q(t)$

• In centralized case:

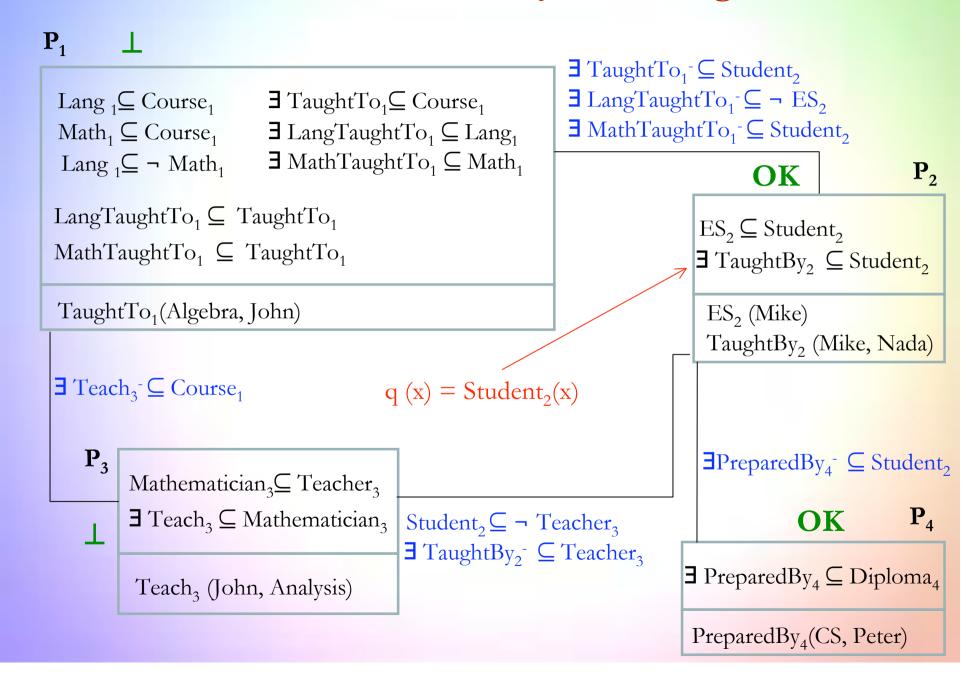
Algorithm PerfectRef: D.Calvenese and al. (JAR 07)

"Tractable reasoning and efficient query answering in description logics"

Data consistency checking

Query Answering

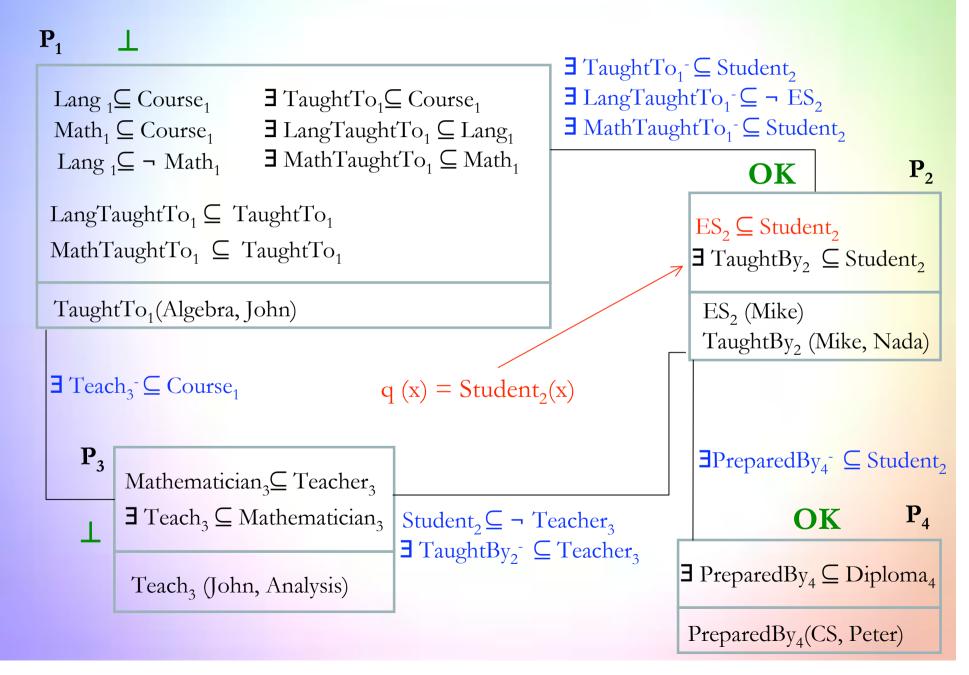
- Query Reformulation
- Reformulation Evaluation



- Contribution:
 - Decentralized algorithm for query answering.
 - Well Founded Answers even in globally inconsistent Networks
 - 1. Decentralized Query Reformulation
 - 2. Checking data consistency w.r.t the peers involved in each reformulation
 - We keep only sound reformulations
 - 3. Evaluation of the sound reformulations
 - Well founded answers

- Contribution:
 - Decentralized algorithm for query answering.
 - ➤ Query reformulation

$$q(x) = Student_2(x)$$



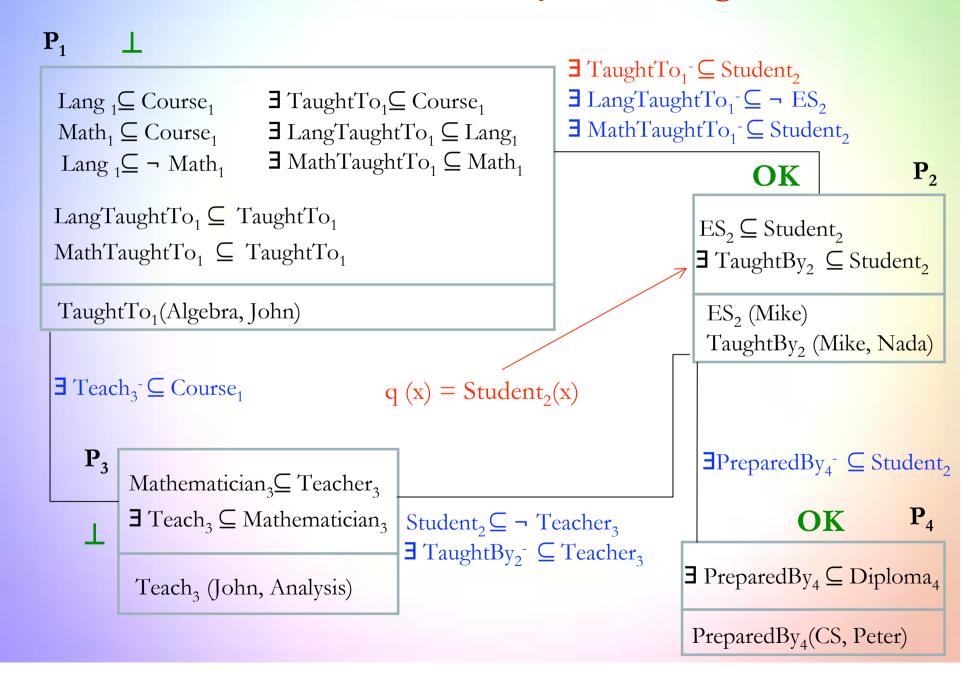
• Contribution:

Decentralized algorithm for query answering.

➤ Query reformulation

$$q(x) = Student_2(x)$$

$$q_1(x) = ES_2(x)$$

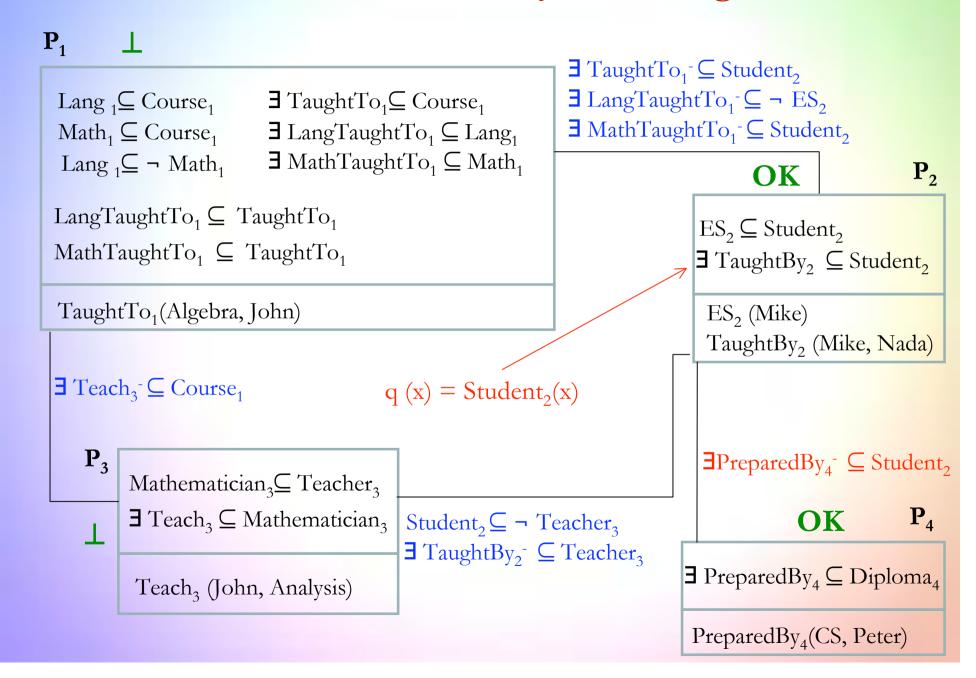


• Contribution:

Decentralized algorithm for query answering.

➤ Query reformulation

```
q(x) = Student_2(x)
q_1(x) = ES_2(x)
q_2(x) = \exists z TaughtTo_1(z, x)
```



• Contribution:

Decentralized algorithm for query answering.

➤ Query reformulation

```
q(x) = Student_2(x)
q_1(x) = ES_2(x)
q_2(x) = \exists z TaughtTo_1(z, x)
q_3(x) = \exists y PreparedBy_4(y, x)
```

• Contribution:

Decentralized algorithm for query answering.

➤ Query reformulation

```
q(x) = Student_2(x)
q_1(x) = ES_2(x)
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Data consistency checking

• Contribution:

Decentralized algorithm for query answering.

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q(x) = Student_2(x)
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Data consistency checking

Consistent algorithm on P₁, P₂ and P₄

On P_1 : \perp

On P₂: OK

On P₄: OK

• Contribution:

Decentralized algorithm for query answering.

➤ Query reformulation

$$q(x) = Student_2(x)$$

$$q_1(x) = ES_2(x)$$

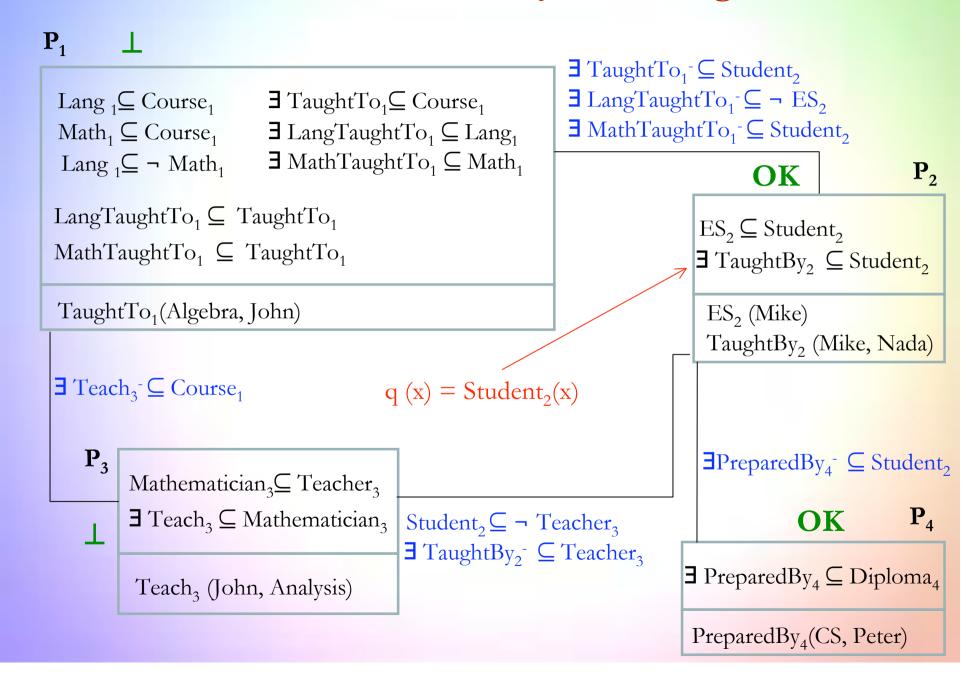
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Data consistency checking

Consistent algorithm on P₁, P₂ and P₄

On P_1 : \perp On P_2 : OK On P_4 : OK Query Evaluation



• Contribution:

Decentralized algorithm for query answering.

> Query reformulation

$$q(x) = Student_2(x)$$
 \longrightarrow

 $q_1(x) = ES_2(x)$

 $q_2(x) = \exists z \text{ TaughtTo}_1(z, x)$

 $q_3(x) = \exists y \text{ PreparedBy}_4(y, x)$

➤ Data consistency checking

Consistent algorithm on P₁, P₂ and P₄

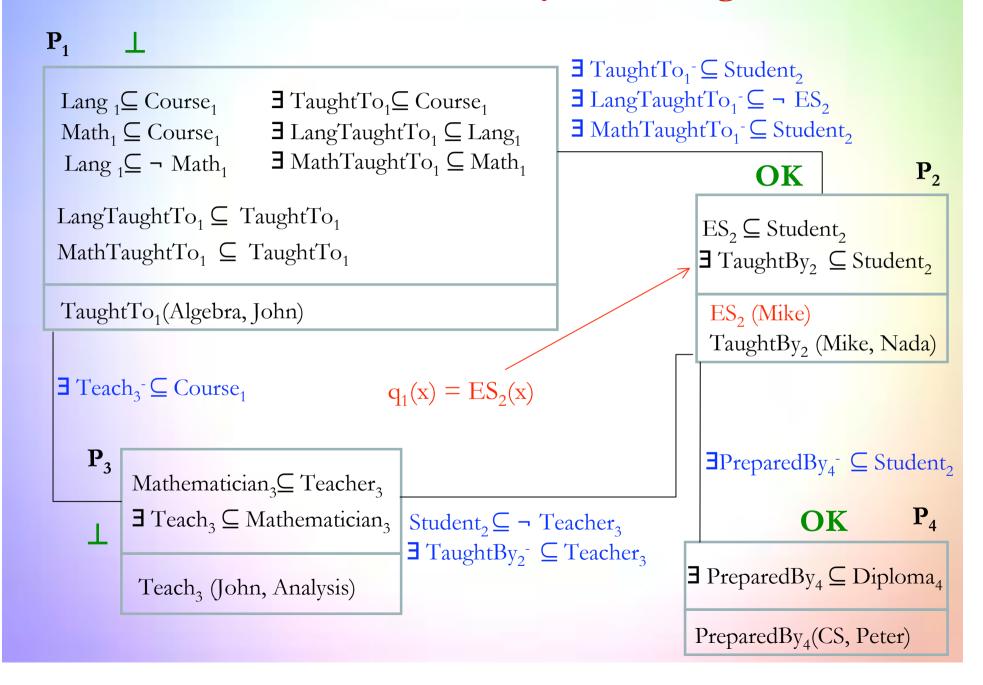
On P_1 : \perp

On P₂: OK

On P₄: OK

Query Evaluation

No answers



• Contribution:

Decentralized algorithm for query answering.

➤ Query reformulation

$$q(x) = Student_2(x)$$

$$q_1(x) = ES_2(x) \qquad \qquad \dots$$

$$q_2(x) = \exists z \text{ TaughtTo}_1(z, x)$$

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Data consistency checking

Consistent algorithm on P₁, P₂ and P₄

On P_1 : \perp

On P₂: OK

On P₄: OK

Query Evaluation

No answers

Well founded Answer= {Mike}

• Contribution:

Decentralized algorithm for query answering.

➤ Query reformulation

$$q(x) = Student_2(x)$$

$$q_1(x) = ES_2(x)$$

$$q_2(x) = \exists z \text{ TaughtTo}_1(z, x)$$

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Data consistency checking

Consistent algorithm on P₁, P₂ and P₄

On P_1 : \perp

On P₂: OK

On P₄: OK

Query Evaluation

No answers

Well founded Answer= {Mike}

```
\mathbf{P}_{1}
                                                                                                     \exists TaughtTo<sub>1</sub> \subseteq Student<sub>2</sub>
                                                                                                    \exists LangTaughtTo<sub>1</sub>-\subseteq \neg ES<sub>2</sub>
                                           ∃ TaughtTo<sub>1</sub>⊆ Course<sub>1</sub>
  Lang _1\subseteq Course
                                                                                                    ∃ MathTaughtTo<sub>1</sub>-⊆ Student<sub>2</sub>
                                           \exists LangTaughtTo<sub>1</sub> \subseteq Lang<sub>1</sub>
  Math_1 \subseteq Course_1
   Lang _1 \subseteq \neg Math<sub>1</sub>
                                           \exists MathTaughtTo<sub>1</sub> \subseteq Math<sub>1</sub>
                                                                                                                                                                      \mathbf{P}_2
                                                                                                                                        OK
  LangTaughtTo<sub>1</sub> \subseteq TaughtTo<sub>1</sub>
                                                                                                                             ES_2 \subseteq Student_2
  MathTaughtTo_1 \subseteq TaughtTo_1
                                                                                                                             \exists TaughtBy<sub>2</sub> \subseteq Student<sub>2</sub>
   TaughtTo<sub>1</sub>(Algebra, John)
                                                                                                                               ES<sub>2</sub> (Mike)
                                                                                                                               TaughtBy, (Mike, Nada)
                                                     q_3(x) = \exists y \text{ PreparedBy}_4(y, x)
  \exists \operatorname{Teach}_{3} \subseteq \operatorname{Course}_{1}
        \mathbf{P}_3
                                                                                                                              \exists Prepared By_4^- \subseteq Student_2
                Mathematician₃⊆ Teacher₃
                                                                                                                                                                      \mathbf{P}_4
                \exists Teach<sub>3</sub> \subseteq Mathematician<sub>3</sub> \mid Student<sub>2</sub> \subseteq \neg Teacher<sub>3</sub>
                                                                                                                                                OK
                                                                     ∃ TaughtBy<sub>2</sub> \subseteq Teacher<sub>3</sub>
                                                                                                                           \exists PreparedBy<sub>4</sub> \subseteq Diploma<sub>4</sub>
                  Teach<sub>3</sub> (John, Analysis)
                                                                                                                            PreparedBy<sub>4</sub>(CS, Peter)
```

• Contribution:

Decentralized algorithm for query answering.

> Query reformulation

$$q(x) = Student_2(x)$$

$$q_1(x) = ES_2(x)$$

$$q_2(x) = \exists z \text{ TaughtTo}_1(z, x)$$

$$q_3(x) = \exists y \text{ PreparedBy}_4(y, x)$$

Data consistency checking

Consistent algorithm on P₁, P₂ and P₄

On P_1 : \perp

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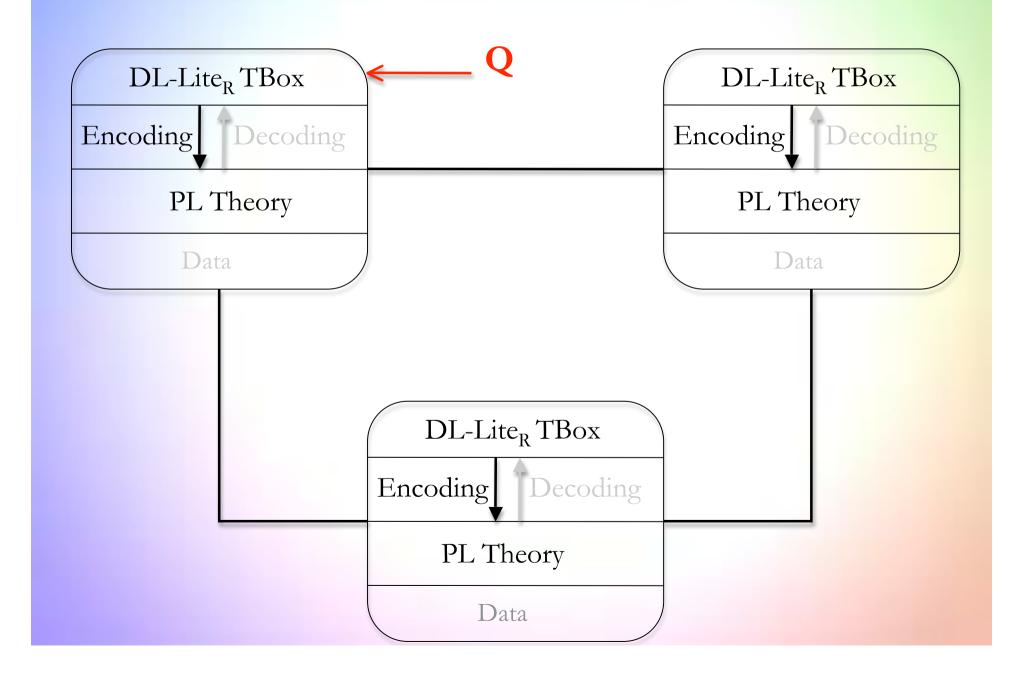
Query Evaluation

No answers

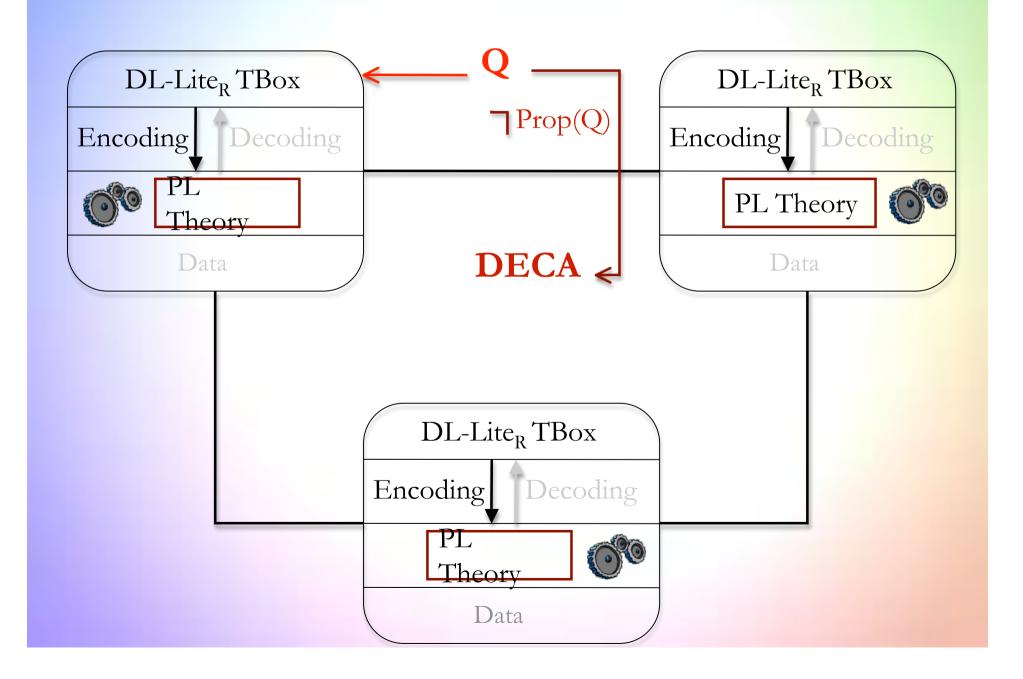
Well founded Answer= {Mike}

Well founded Answer= {Peter}

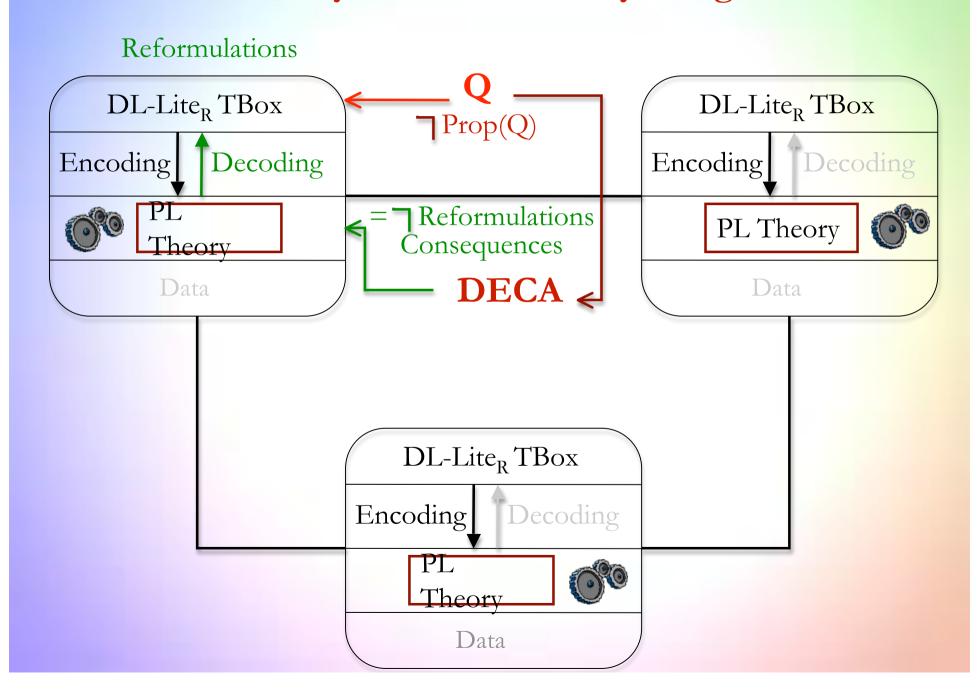
Decentralized Query Reformulation by using SomeWhere



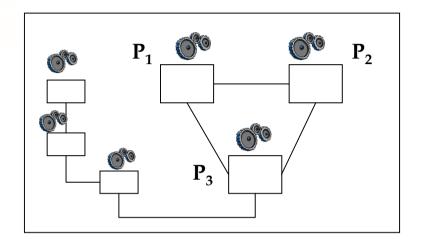
Decentralized Query Reformulation by using SomeWhere



Decentralized Query Reformulation by using SomeWhere



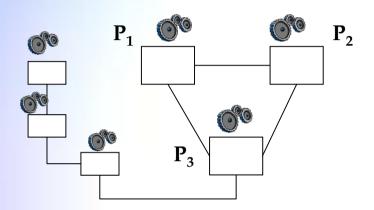
Outline



- ☐ DL-Lite_R decentralized data management system
 - ☐ Data consistency checking
 - ☐ Query Answering
 - ☐ Conclusion & Future Work

Conclusion & future work

1. DL-Lite_R decentralized data model



- 2. Decentralized Algorithms for:
 - ✓ Data consistency checking
 - **✓** Query Answering
- 3. Decentralized & Centralized Algorithms for:
 - ✓ View consistency checking
 - **✓ Query Answering using views**

Other Applications

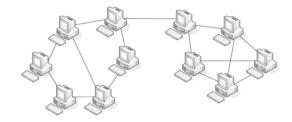
- Mapping discovery
- Data reconciliation
- Explanation of answers to queries

Feasibility of the approach with other dialects of DL-Lite

Thank You

Questions?

SomeWhere and the Semantic Web



After SomeOWL and SomeRDFS

SomeDL-Lite_R

Nada Abdallah